

INTERNET PROTOCOL (IP)

Packet-Level Functionality and Fields of the Internet Protocol

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IP Functionality:

IP offers end-to-end connectivity between devices on separate networks. An IP header enables a packet to be routed through an internetwork using a software address that includes a network address portion and a host address portion. Routers along a path examine the destination IP address found in the IP header to determine if they can forward the packet.

The IP header also includes fields to ensure the packet does not circulate endlessly in case of an internetwork loop and fields to enable a packet to be broken down (fragmented) into smaller pieces when the network path does not support a single MTU (maximum transmission unit) size from end-to-end.

The currently widely-implemented version of IP is IPv4. IPv6 is in development. (See www.packet-level.com for more information.)

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NOTES:

Note #1 Version field: This 4-bit field indicates the version of IP in use. The following lists the version numbers assigned:

v0 Reserved	V7 TP/DC: The Next Internet
V4 Internet Protocol (RFC 791)	v0 PIP: The P Internet Protocol
v6 ST (Stream) Datagram Mode (RFC 1190)	v9 TUBA
v0 Internet Protocol Version 6 (RFC 1752)	V15 Reserved

Versions 1-3 and 10-14 are unassigned.

Note #2 Internet Header Length (IHL) field: This field denotes the length of the IP header in 4-byte increments. For example, the value 5 in this field defines a 20-byte IP header. Subtract this value from the Total Length field to obtain the length of transport layer and data in the packet.

Note #3 Precedence/Type of Service field: This field identifies the quality of service that a packet should receive (if possible). Bits 0-2 are used for Precedence; bits 3-6 are used for Type of Service; bit 7 is reserved.

Precedence	Description	TOS Value	Service Description
0000	Default (no specific service type)	0000	
0001	Network Control	0001	Minimize Monetary Cost
0010	Internetwork Control	0010	Maximize Reliability
0101	CRITICAL/ECF	0100	Maximize Throughput
1000	Flash Override	1000	Minimize Delay
0111	Flash	1111	Maximize Security
0110	Immediate		
0011	Priority		

Note #4 Total Length field: This field defines the total number of bytes in the IP packet from the start of the IP header through the end of valid data. This field does not include any bytes inserted as padding at the end of the packet.

Note #5 Flags field: This field actually consists of three separate single-bit fields. The fields are:

Bit 0: Reserved (set to 0)
Bit 1: May/Don't Fragment bit (0=May Fragment; 1=Don't Fragment)
Bit 2: More/Last Fragment bit (0=Last Fragment; 1=More Fragments to Come)

The following figure denotes the bit placement/value as well.

Bit 0	Bit 1	Bit 2
0	D	M
0	F	F

DF = Don't Fragment bit
MF = More Fragment bit

Note #6 Time to Live field: This field defines the remaining lifetime of the packet. Although originally defined as a measurement of "seconds" it is often referred to as the "remaining hop count" value since routing devices will decrement this field value by 1 even if it didn't take one second for the routing process to transpire. Typical TTL starting values are 32, 64 and 128.

A packet with a TTL value of 1 arriving at a router cannot be routed (the router may not decrement the TTL value to 0 and forward the packet). A packet with a TTL value of 1 arriving at the destination can be processed, however.

In the case of a fragmented packet, all fragments are given the TTL value at the time of fragmentation. The fragments may take different paths to the destination and may end up with different TTL values upon arrival. All fragments must arrive at the destination within the TTL value of the first-received fragment, however.

Note #7 Header Checksum field: This field contains a checksum value for the IP header only and includes all fields except the Header Checksum field itself.

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Class-based IP Addressing

Class-based addressing offers an IP address system that separates address into several basic classes with predefined field lengths for the network portion of the address and the host ID portion of the address as shown below.

Class	Net/Host	Network Mask	First Byte in Binary	First Byte in Decimal
Class A	N.H.H.H	255.0.0.0	0xxxxxxx	0 - 127*
Class B	N.N.H.H	255.255.0.0	10xxxxxx	128 - 191
Class C	N.N.N.H	255.255.255.0	110xxxxx	192 - 223
Class D	Multicast	n/a	1110xxxx	224 - 239
Class E	Reserved	n/a	1111xxxx	240 - 255

* A network address starting with 0 is invalid. The network address 127 is reserved for the localhost address.

When a single network address is required to support two or more physical networks, the address can be subnetted. Subnetting is covered in detail in RFC 1918.

The "Network Address"

Placing the value 0 (binary) in all positions of the host ID portion of the network's address. For example, 9.0.0.0 can be used to refer to the entire 9.x.x.x network. It is not allowable to assign a host the host ID of all 0s.

All-Network Broadcast

Placing the value 1 (binary) in all positions of the network and host ID portion of an address is a network broadcast. The decimal value is 255.255.255.255. Broadcasts are typically not be forwarded by routers (even though the broadcast is directed to "all networks").

Subnet Broadcast

Placing the value 1 (binary) in all positions of the host ID portion of an address is a broadcast for a specific subnetwork. For example, 9.255.255.255 is a broadcast onto network 9.0.0.0.

Classless IP Addressing

Classless Interdomain Routing (CIDR) defines a method for addressing without Class A, B, or C network delineations. Instead of using such rigid delimiters and making assumptions on the bytes used for network and host portion of the address, classless addresses consist of an address and the prefix. The prefix indicates the number of bits to be masked off for the network portion. The following table illustrates the classless addressing prefixes and masks.

CIDR Prefix	Decimal Mask	Binary
/0	0.0.0.0	00000000 00000000 00000000 00000000
/1	128.0.0.0	10000000 00000000 00000000 00000000
/2	192.0.0.0	11000000 00000000 00000000 00000000
/3	224.0.0.0	11100000 00000000 00000000 00000000
/4	240.0.0.0	11110000 00000000 00000000 00000000
/5	248.0.0.0	11111000 00000000 00000000 00000000
/6	252.0.0.0	11111100 00000000 00000000 00000000
/7	254.0.0.0	11111110 00000000 00000000 00000000
/8	255.0.0.0	11111111 00000000 00000000 00000000
/9	255.128.0.0	11111111 10000000 00000000 00000000
/10	255.192.0.0	11111111 11000000 00000000 00000000
/11	255.224.0.0	11111111 11100000 00000000 00000000

(and so on...)

Private Addresses

The Internet Assigned Numbers Authority (IANA) has reserved the following addresses for private internets.

10.0.0.0 - 10.255.255.255 (10/8 prefix)
172.16.0.0 - 172.31.255.255 (172/16/12 prefix)
192.168.0.0 - 192.168.255.255 (192/16/16 prefix)

These numbers cannot be advertised on the Internet. Network Address Translation (NAT) systems may be used to allow connectivity between private from public addressing systems.

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Multicast Addressing

Multicasting is used to address a group of IP devices. The multicast address can only be used in the destination IP address field (never in the source IP address field). Routing Information Protocol (RIP) version 2, Open Shortest Path First (OSPF), Service Location Protocol (SLP) and Internet Group Management Protocol (IGMP) use multicast addresses. See RFC 1112, Host Extensions for IP Multicasting* for more information.

All IP multicast addresses are assigned between 224.0.0.0 and 239.255.255.255. The range 224.0.0.0 through 224.0.0.255, inclusive, are reserved for the use of routing protocols and other topology and maintenance discovery protocols. These multicast packets should not be forwarded by multicast routers.

Multicast MAC Headers

The IP multicast address provides the "seed" for the creation of a multicast MAC address by translating the last three bytes of the IP multicast address into hexadecimal and placed in the last three bytes of the MAC address.

For example, in the packet shown at right, an IGMP packet addressed to IP address 224.0.1.22 (Service Location Protocol multicast address) uses the MAC header destination address 0x000000000016 (where '16' is the hexadecimal equivalent of '22' decimal). The leading three bytes use an IANA-assigned MAC identifier (0x00000000) with the first byte changed to an odd number - 0x01 - to denote a multicast MAC packet).

Partial List of Assigned Multicast Addresses

The following list includes the basic set of multicast addresses (from 224.0.0.0 to 224.0.0.255). Refer to www.iana.org for a complete list of assigned multicast addresses.

Address	Description	Address	Description (continued)
224.0.0.0	Base Address (Reserved)	224.0.0.13	All PIM Routers
224.0.0.1	All Systems on this Subnet	224.0.0.14	RSVP/ENCAPSULATION
224.0.0.2	All Routers on this Subnet	224.0.0.15	All-OSPF-routers
224.0.0.3	Unassigned	224.0.0.16	designated-ebm
224.0.0.4	DVMRP Routers	224.0.0.17	all-ebms
224.0.0.5	OSPF/IGP All Routers	224.0.0.18	VRRP
224.0.0.6	OSPF/IGP Designated Rtr.	224.0.0.19	IPALL/15e
224.0.0.7	ST Routers	224.0.0.20	IPALL/21e
224.0.0.8	ST Hoses	224.0.0.21	IPAll Intermediate Systems
224.0.0.9	RIP2 Routers	224.0.0.22	IGMP
224.0.0.10	IGRP Routers	224.0.0.23	224.0.0.23-224.0.0.255 unassigned
224.0.0.11	Mobile-Agent		

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Multicast Packet (Joining SLP Group)

DLC: Destination = Multicast 01005E000116
DLC: Source = Station 00A0C3C0C8B
DLC: Ethertype = 0800 (IP)
IP: Version = 4, header length = 24 bytes
IP: Type of service = 00
IP: Total length = 32 bytes
IP: Identification = 768
IP: Flags = 0X
IP: .0. = may fragment
IP: .0. = last fragment
IP: Fragment offset = 0 bytes
IP: Time to live = 1 seconds/hops
IP: Protocol = 2 (IGMP)
IP: Header checksum = 365F (correct)
IP: Source address = [10.0.0.99]
IP: Destination address = [224.0.1.22]
IP: Options follow
IP: Option 148
IP: 13 byte(s) of header padding
IGMP: Version = 1
IGMP: Type = 6 (Ver2 Membership Report)
IGMP: Unused = 0x00
IGMP: Checksum = 0B89 (correct)
IGMP: Group Address = [224.0.1.22]

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Basic TCP/IP Stack Elements

The diagram shows the layers of the TCP/IP stack:

- Upper-Layer Protocols (such as SNMP, HTTP, FTP, telnet, etc.)
- User Datagram Protocol (UDP) RFC 768 (connectionless transport)
- Transmission Control Protocol (TCP) RFC 793 (connection-oriented transport)
- Internet Control Messaging Protocol (ICMP) RFC 792
- Internet Protocol (IP) RFC 791
- Address Resolution Protocol (ARP) RFC 826
- Media Access Control (MAC Layer) (Ethernet, Token Ring, FDDI, etc.)

UDP-over-IP Communication (DNS Query)

DLC: Destination = Station 00104B30C44A
DLC: Source = Station 00A0C3C0C8B
DLC: Ethertype = 0800 (IP)
IP: Vers. = 4, header length = 20 bytes
IP: Type of service = 00
IP: Total length = 58 bytes
IP: Identification = 45568
IP: Flags = 0X may/last fragment
IP: Fragment offset = 0 bytes
IP: Time to live = 128 seconds/hops
IP: Protocol = 17 (UDP)
IP: Header checksum = 744F (correct)
IP: Source address = [10.0.0.99]
IP: Destination address = [10.0.0.1]
IP: No options
UDP: Source port = 1031
UDP: Destination port = 53 (Domain)
UDP: Length = 38
UDP: Checksum = CB03 (correct)
UDP: [30 byte(s) of data]
DNS: DNS: ID = 1
DNS: Flags = 01 Query: Recursion desired
DNS: Flags = 0X; verified data only
DNS: Question count = 1, Answer count = 0
DNS: Authority count = 0, Addl. records = 0
DNS: ZONE Section
DNS: Name = ftppool1.NAI
DNS: Type = Host address (A,1)
DNS: Class = Internet (IN,1)
DNS:

TCP-over-IP Communication (TCP Handshake)

DLC: Destination = Station Trndnt3C3D8B
DLC: Source = Station Eagle 005FC1
DLC: Ethertype = 0800 (IP)
IP: Version = 4, header length = 20 bytes
IP: Type of service = 00
IP: Total length = 44 bytes
IP: Identification = 1536
IP: Flags = 4X; don't/last fragment
IP: Fragment offset = 0 bytes SV
IP: Time to live = 32 seconds/hops
IP: Protocol = 6 (TCP)
IP: Header checksum = 284F (correct)
IP: Source address = [130.57.20.10]
IP: Destination address = [130.57.20.1]
IP: No options
TCP: Source port = 1026
TCP: Destination port = 524
TCP: Initial sequence number = 12952
TCP: Next expected Seq number = 12953
TCP: Data offset = 24 bytes
TCP: Flags =
...0. = (No urgent pointer)
...0. = (No acknowledgment)
...0. = (No push)
...0. = (No reset)
...1. = (No FIN)
TCP: Window = 8192
TCP: Checksum = 1303 (correct)
TCP: Options follow
TCP: Maximum segment size = 1460

IP Header Structure

The diagram shows the structure of the IP header:

- Version (4 bits) - See Note #1
- IHL (4 bits) - See Note #2
- Precedence/Type of Service (8 bits) - See Note #3
- Total Length (16 bits) - See Note #4
- Identification (16 bits) - See 'IP Fragmentation' Sidebar
- Flags (3 bits) - See Note #5
- Fragment Offset (13 bits) - See 'IP Fragmentation' Sidebar
- Time to Live (8 bits) - See Note #6
- Protocol (8 bits) - See Table #1
- Header Checksum (16 bits) - See Note #7
- Source IP Address (32 bits) - See 'Class-based/Classless IP Addressing' Sidebar
- Destination IP Address (32 bits) - See 'Class-based/Classless IP Addressing' and Multicast Addressing' Sidebars
- Options (if any) (variable length) - See Table #2
- Data (variable length)

Table #1: Protocol Field

Decimal	Protocol	Decimal	Protocol (continued)
0	IPv6 Hop-by-Hop Option	36	XTP
1	Internet Control Message Protocol	37	Datagram Delivery Protocol
2	Internet Group Management Protocol	38	IDRP Control Message Transport
3	Gateway-to-Gateway	39	TP++ Transport Protocol
4	IP in IP (encapsulation)	40	IL Transport Protocol
5	Stream	41	IPv6
6	Transmission Control Protocol	42	Source Demand Routing Protocol
7	CBT	43	IPv6-Route Routing Header for IPv6
8	Exterior Gateway Protocol	44	Fragment Header for IPv6
9	Any private interior gateway (i.e. IGRP)	45	Inter-Domain Routing Protocol
10	BBN RZC Monitoring	46	Reservation Protocol
11	Network Voice Protocol	47	General Routing Encapsulation
12	PUP	48	Mobile Host Routing Protocol
13	ARGUS	49	BNA
14	EMCON	50	Encap Security Payload for IPv6
15	Cross Net Debugger	51	Authentication Header for IPv6
16	Chaos	52	Integrated Net Layer Security
17	User Datagram Protocol	53	IP with Encryption
18	Multiplexing	54	NSMA Address Resolution Protocol
19	DON Measurement Subsystem	55	IP Mobility
20	Host Monitoring	56	Transport Layer Security Protocol
21	Packet Radio Measurement	57	SKIP
22	XEROX NS IDP	58	ICMP for IPv6
23	Trunk-1	59	No Next Header for IPv6
24	Trunk-2	60	Destination Options for IPv6
25	Leaf-1	61	Any host internal protocol
26	Leaf-2	62	CTFP
27	Reliable Data Protocol	63	Any local network
28	Internet Reliable Transaction	64	SATNET and Backroom EXPAK
29	ISO Transport Protocol Class 4	65	Krypsolam
30	Bulk Data Transfer Protocol	66	MIT Remote Virtual Disk Protocol
31	MFE Network Services Protocol	67	Internet Pluribus Packet Core
32	MERIT Internal Protocol	68	Any distributed file system
33	Sequential Exchange Protocol	69	SATNET Monitoring
34	Third Party Connect Protocol	70	VISA Protocol
35	Inter-Domain Policy Routing Protocol		
36	XTP		

Table #2: IP Header Options

#	Value	Reference	#	Value	Reference (continued)
0	0	End of Options List	13	205	Experimental Flow Control
1	1	No Operation	14	142	Experimental Assoc Control
2	130	Security	15	15	Encode
3	131	Loose Source Route	16	144	IMI Traffic Descriptor
4	68	Time Stamp	17	145	Extended Internet Protocol
5	133	Extended Security	18	82	Traceroute
6	134	Commercial Security	19	147	Address Extension
7	7	Record Route	20	148	Router Alert
8	136	Stream ID	21	149	Selective Directed Broadcast
9	137	Strict Source Route	22	150	NSAP Address
10	10	Experimental Measurement	23	151	Dynamic Packet State
11	11	MTU Probe	24	152	Upstream Multicast Pkt.
12	12	MTU Reply			

IP Fragmentation

Fragmentation is necessary when a packet is too large to fit on the media. The packet must be fragmented into several pieces at the point where the media size changes (typically a router) and reassembled at the destination.

For example, consider a device that sends a 4096-byte packet from a Token Ring network to a device located one router on an Ethernet network. The Ethernet network can only support 1518 byte packets. The router that connects the Token Ring network to the Ethernet network must fragment a single packet into three separate packets.

The process of fragmentation uses three fields of the IP header: the Identification field, Flags field, and the Fragment Offset field.

Identification field: All IP packets contain a unique ID number. When a packet must be fragmented into several pieces, the same ID number is placed in each fragment to identify them as a set.

Flags field: This field consists of three bits - only two are used. Bits 1 and 2 of this field (see Note #5) are used to indicate if fragmentation is allowed (Don't Fragment bit) and if a packet is part of a fragment set (More Fragments).

Fragmentation Offset field: This field indicates where the fragment's data should be placed in relation to the other fragments data. The field indicates the offset in a units of 8 octets (64 bits).

Fragment 1 of 3

IP: ----- IP Header -----
IP: Total length = 1500 bytes
IP: Identification = 19485
IP: Flags = 2X
IP: .0. = may fragment
IP: .1. = more fragments
IP: Fragment offset = 0 bytes
IP: Time to live = 32 seconds/hops
IP: Protocol = 1 (ICMP)
IP: Multi-Frame IP data, Frames: 1, 2, 3

Fragment 2 of 3

IP: ----- IP Header -----
IP: Total length = 1500 bytes
IP: Identification = 19485
IP: Flags = 2X
IP: .0. = may fragment
IP: .1. = more fragments
IP: Fragment offset = 1480 bytes
IP: Time to live = 32 seconds/hops
IP: 1480 bytes of data continuation of IP ident=19485

Fragment 3 of 3

IP: ----- IP Header -----
IP: Total length = 1164 bytes
IP: Identification = 19485
IP: Flags = 0X
IP: .0. = may fragment
IP: .0. = last fragment
IP: Fragment offset = 2960 bytes
IP: Time to live = 32 seconds/hops
IP: [1144 bytes of data continuation of IP ident=19485]

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